

Minimal Deductive Systems for RDF

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RDF Language and Model Theory looks a little bit complicated at first sight...



RDF Semantics

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RDF Semantics

W3C Recommendation 10

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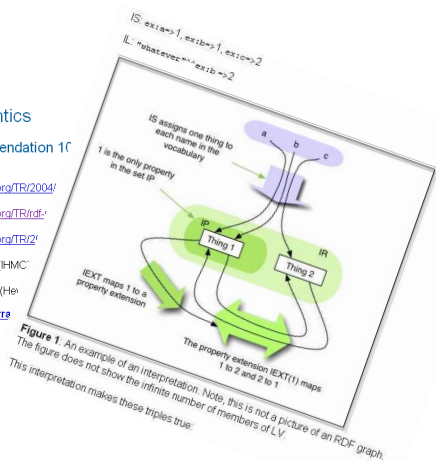
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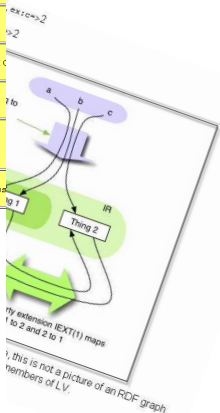


RDF Language and Model Theory looks a little bit complicated at first sight...

```
EXT( $\{(rdfs:subPropertyOf)\}$ ) is transitive and reflexive on IP  
If  $\langle x,y \rangle$  is in EXT( $\{(rdfs:subPropertyOf)\}$ ) then  $x$  and  $y$  are in IP and EXT( $x$ ) is a subset  
If  $x$  is in IC then  $\langle x, \{(rdfs:Resource)\} \rangle$  is in EXT( $\{(rdfs:subClassOf)\}$ )  
If  $\langle x,y \rangle$  is in EXT( $\{(rdfs:subClassOf)\}$ ) then  $x$  and  $y$  are in IC and ICEXT( $x$ ) is a subset  
EXT( $\{(rdfs:subClassOf)\}$ ) is transitive and reflexive on IC  
If  $x$  is in ICEXT( $\{(rdfs:ContainerMembershipProperty)\}$ ) then  
 $\langle x, \{(rdfs:member)\} \rangle$  is in EXT( $\{(rdfs:subPropertyOf)\}$ )  
If  $x$  is in ICEXT( $\{(rdfs:Datatype)\}$ ) then  $\langle x, \{(rdfs:Literal)\} \rangle$  is in EXT( $\{(rdfs:subClassOf)\}$ )
```

RDFS axiomatic triples.

```
rdf:type rdfs:domain rdfs:Resource .  
rdfs:domain rdfs:domain rdf:Property .  
rdfs:range rdfs:domain rdf:Property .  
rdfs:subPropertyOf rdfs:domain rdf:Property .  
rdfs:subClassOf rdfs:domain rdfs:Class .  
rdf:subject rdfs:domain rdf:Statement .  
rdf:predicate rdfs:domain rdf:Statement .  
rdf:object rdfs:domain rdf:Statement .  
rdfs:member rdfs:domain rdfs:Resource .  
rdf:first rdfs:domain rdf:List .  
rdf:rest rdfs:domain rdf:List .  
rdfs:seeAlso rdfs:domain rdfs:Resource .  
rdfs:isDefinedBy rdfs:domain rdfs:Resource .  
rdfs:comment rdfs:domain rdfs:Resource .
```



RDF Language and Model Theory looks a little bit complicated at first sight...

`!EXT((rdf:subPropertyOf))` is transitive and reflexive on IP

If $\langle x,y \rangle$ is in `!EXT((rdf:subPropertyOf))` then x and y are in IP and `!EXT((rdf:subPropertyOf))` is a subset of `!IP`

If x is in IC then $\langle x, !((rdf:Resource)) \rangle$ is in `!EXT((rdf:subPropertyOf))`

If $\langle x,y \rangle$ is in `!EXT((rdf:subPropertyOf))` then x and y are in IP and `!EXT((rdf:subPropertyOf))` is a subset of `!IP`

`!EXT((rdf:subPropertyOf))` is transitive and reflexive on IP

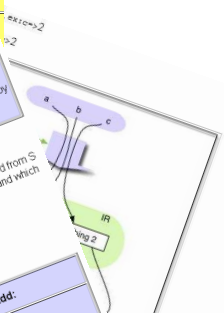
RDF entailment rules	
Rule Name	then add
rdf1	<code>aaa rdf:type rdf:Property .</code>
rdf2	<code>...rnn rdf:type rdf:XMLLiteral .</code> where <code>...rnn</code> identifies a blank node allocated to <code>ll</code> by rule <code>lg</code> .

These rules are complete in the following sense.
RDF entailment lemma `S` `rdf`-entails `E` if and only if there is a graph which can be derived from `S` plus the RDF axiomatic triples by the application of rule `lg` and the RDF entailment rules and which simply entails `E`. (Proof in Appendix A)

Note that this does not require the use of rule `gl`.

7.3 RDFS Entailment Rules

RDFS entailment rules.	
Rule Name	then add:
rdfs1	<code>...rnn rdf:type rdfs:Literal .</code> where <code>...rnn</code> identifies a blank node allocated to <code>ll</code> by rule <code>lg</code> .



RDF Language and Model Theory looks a little bit complicated at first sight...

```

IEXT((rdf:subPropertyOf)) is transitive and reflexive on IP
if <x,y> is in IEXT((rdf:subPropertyOf)) then x and y are in IP and IEXT((rdf:subPropertyOf)) is transitive and reflexive on IP
if x is in IC
if <x,y> is in IEXT((rdf:subPropertyOf)) then x and y are in IP and IEXT((rdf:subPropertyOf)) is transitive and reflexive on IP
IEXT((rdf:subPropertyOf)) is transitive and reflexive on IP

```

RDF entailment rules

Rule Name	E cont
rdf1	uuu aaa
rdf2	uuu

These rules plus the simply Note

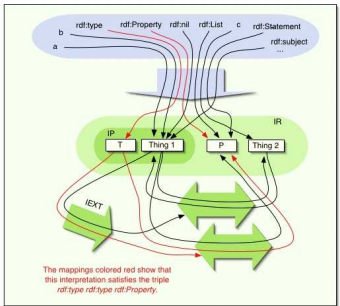


Figure 2: An rdf-interpretation.

This is not the smallest rdf-interpretation which extends the earlier example, since one made IEXT(T) be {<1,2>, <T,2>}, and managed without having P in the universe. In gettitude of an RDF graph.

RDF's entailment

Rule Name	uuu aaa ll.
rdf1	where ll is a plain literal (with or without a language tag).



Several dozens of reserved keywords...

rdfs:Resource [res]	rdf:type [type]	rdfs:isDefinedBy [isDefined]
rdf:Property [prop]	rdfs:domain [dom]	rdfs:comment [comment]
rdfs:Class [class]	rdfs:range [range]	rdfs:label [label]
rdfs:Literal [literal]	rdfs:subClassOf [sc]	rdf:value [value]
rdfs:Datatype [datatype]	rdfs:subPropertyOf [sp]	rdf:nil [nil]
rdf:XMLLiteral [xmlLit]	rdf:subject [subj]	rdf:_1 [_1]
rdfs:Container [cont]	rdf:predicate [pred]	rdf:_2 [_2]
rdf:Statement [stat]	rdf:object [obj]	...
rdf:List [list]	rdfs:member [member]	rdf:_i [_i]
rdf:Alt [alt]	rdf:first [first]	...
rdf:Bag [bag]	rdf:rest [rest]	
rdf:Seq [seq]	rdfs:seeAlso [seeAlso]	
rdfs:ContainerMembershipProperty [contMP]		

...plus axiomatic triples...

(type,	type,	prop)
(subj,	type,	prop)
(pred,	type,	prop)
(obj,	type,	prop)
(first,	type,	prop)
(rest,	type,	prop)
(value,	type,	prop)
(_1,	type,	prop)
(_1,	type,	contMP)
(_2,	type,	prop)
(_2,	type,	contMP)
...		
(_i,	type,	prop)
(_i,	type,	contMP)
...		
(nil,	type,	prop)
(xmlLit,	type,	datatype)

...plus more axiomatic triples...

(type,	dom,	res)	(type,	range,	class)
(dom,	dom,	prop)	(dom,	range,	class)
(range,	dom,	prop)	(range,	range,	class)
(sp,	dom,	prop)	(sp,	range,	prop)
(sc,	dom,	class)	(sc,	range,	class)
(subj,	dom,	stat)	(subj,	range,	res)
(pred,	dom,	stat)	(pred,	range,	res)
(obj,	dom,	stat)	(obj,	range,	res)
(member,	dom,	res)	(member,	range,	res)
(first,	dom,	list)	(first,	range,	res)
(rest,	dom,	list)	(rest,	range,	list)
(seeAlso,	dom,	res)	(seeAlso,	range,	res)
(isDefined,	dom,	res)	(isDefined,	range,	res)
(comment,	dom,	res)	(comment,	range,	literal)
(label,	dom,	res)	(label,	range,	literal)
(value,	dom,	res)	(value,	range,	res)
(_1,	dom,	res)	(_1,	range,	res)
(_2,	dom,	res)	(_2,	range,	res)
...			...		
(_i,	dom,	res)	(_i,	range,	res)
...			...		

...plus more axiomatic triples...

```
(alt,      sc,      cont)
(bag,      sc,      cont)
(seq,      sc,      cont)
(contMP,   sc,      prop)
(xmlLit,   sc,      literal)
(datatype, sc,      class)

(isDefined, sp,     seeAlso)
```

...and on top of this a (slightly) non-standard model theory

- ▶ A notion of **interpretation**

(Res, Prop, Class, Lit, PExt, CExt, Int)

including subsets of the universe denoting properties, classes and literals, and mapping defining their extensions.

- ▶ Notion of interpretation of blank nodes
- ▶ Definition of reflexivity, transitivity and semi-extensionality of subClass and subProperty
- ▶ Typing restrictions

But...we need a workable language to bring to reality the vision of the Semantic Web

Would like to:

- ▶ Have a **simple user-language** to be able to popularize RDF among Web users.
- ▶ Have a **simple specification** to allow sound development work.
- ▶ Have a **language in streamlined form** to make it easy to formalize and prove results about its properties.

What is to be done?: To simplify the language

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There is a minimal fragment of the theory
preserving the essential core of RDFS

What is to be done?: To simplify the language



There is a minimal fragment of the theory preserving the essential core of RDFS

- ▶ Basic idea: Separate user-language from features and constructors which define and specify the language.
- ▶ Concentrate in vocabulary with non-trivial semantics.

Main contributions & Outline

- ▶ Identify a fragment of RDFS that covers the crucial vocabulary and preserves the original RDFS semantics.
- ▶ Study dependencies among vocabulary and develop sound and complete proof systems for each fragment.
- ▶ Present algorithms to modularize reasoning according to relevant vocabulary.

Main contributions & Outline

- ▶ Identify a fragment of RDFS that covers the crucial vocabulary and preserves the original RDFS semantics.
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Main contributions & Outline

- ▶ Identify a fragment of RDFS that covers the crucial vocabulary and preserves the original RDFS semantics.
- ▶ Study dependencies among vocabulary and develop sound and complete proof systems for each fragment.
- ▶ Present algorithms to modularize reasoning according to relevant vocabulary.

The core vocabulary

<code>rdfs:Resource</code> [<code>res</code>]	<code>rdf:type</code> [<code>type</code>]	<code>rdfs:isDefinedBy</code> [<code>isDefined</code>]
<code>rdf:Property</code> [<code>prop</code>]	<code>rdfs:domain</code> [<code>dom</code>]	<code>rdfs:comment</code> [<code>comment</code>]
<code>rdfs:Class</code> [<code>class</code>]	<code>rdfs:range</code> [<code>range</code>]	<code>rdfs:label</code> [<code>label</code>]
<code>rdfs:Literal</code> [<code>literal</code>]	<code>rdfs:subClassOf</code> [<code>sc</code>]	<code>rdf:value</code> [<code>value</code>]
<code>rdfs:Datatype</code> [<code>datatype</code>]	<code>rdfs:subPropertyOf</code> [<code>sp</code>]	<code>rdf:nil</code> [<code>nil</code>]
<code>rdf:XMLLiteral</code> [<code>xmlLit</code>]	<code>rdf:subject</code> [<code>subj</code>]	<code>rdf:_1</code> [<code>_1</code>]
<code>rdfs:Container</code> [<code>cont</code>]	<code>rdf:predicate</code> [<code>pred</code>]	<code>rdf:_2</code> [<code>_2</code>]
<code>rdf:Statement</code> [<code>stat</code>]	<code>rdf:object</code> [<code>obj</code>]	...
<code>rdf:List</code> [<code>list</code>]	<code>rdfs:member</code> [<code>member</code>]	<code>rdf:_i</code> [<code>_i</code>]
<code>rdf:Alt</code> [<code>alt</code>]	<code>rdf:first</code> [<code>first</code>]	...
<code>rdf:Bag</code> [<code>bag</code>]	<code>rdf:rest</code> [<code>rest</code>]	
<code>rdf:Seq</code> [<code>seq</code>]	<code>rdfs:seeAlso</code> [<code>seeAlso</code>]	
<code>rdfs:ContainerMembershipProperty</code> [<code>contMP</code>]		

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<code>rdfs:Container [cont]</code>	<code>rdf:predicate [pred]</code>	<code>rdf:_2 [_2]</code>
<code>rdf:Statement [stat]</code>	<code>rdf:object [obj]</code>	...
<code>rdf:List [list]</code>	<code>rdfs:member [member]</code>	<code>rdf:_i [_i]</code>
<code>rdf:Alt [alt]</code>	<code>rdf:first [first]</code>	...
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$$\rho_{df} = \{sp, sc, type, dom, range\}$$

Blank node rule

$\frac{G}{H}$ if there is a homomorphism $\mu : H \rightarrow G$

Core rules

Subproperty (transitivity, definition)

$$\frac{(A, \text{sp}, B) (B, \text{sp}, C)}{(A, \text{sp}, C)}$$

$$\frac{(A, \text{sp}, B) (X, A, Y)}{(X, B, Y)}$$

Subclass (transitivity, definition)

$$\frac{(A, \text{sc}, B) (B, \text{sc}, C)}{(A, \text{sc}, C)}$$

$$\frac{(A, \text{sc}, B) (X, \text{type}, A)}{(X, \text{type}, B)}$$

Typing (domain, range)

$$\frac{(A, \text{dom}, B) (X, A, Y)}{(X, \text{type}, B)}$$

$$\frac{(A, \text{range}, B) (X, A, Y)}{(Y, \text{type}, B)}$$

Implicit Typing (strange case...)

$$\frac{(A, \text{dom}, B) (C, \text{sp}, A) (X, C, Y)}{(X, \text{type}, B)}$$

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Reflexivity rules

If “ $(A, \text{type}, \text{Property})$ ” then (A, sp, A)

Subproperty Reflexivity

$$\frac{(X, A, Y)}{(A, \text{sp}, A)}$$

$$\frac{(A, \text{sp}, B)}{(A, \text{sp}, A) (B, \text{sp}, B)}$$

$$\overline{(p, \text{sp}, p)} \text{ for } p \in \rho \text{df}$$

$$\frac{(A, \text{dom}, X)}{(A, \text{sp}, A)} \quad \frac{(A, \text{range}, X)}{(A, \text{sp}, A)}$$

Subclass Reflexivity

$$\frac{(A, \text{sc}, B)}{(A, \text{sc}, A) (B, \text{sc}, B)}$$

$$\frac{(X, \text{dom}, A)}{(A, \text{sc}, A)}$$

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$$\frac{(X, \text{range}, A)}{(A, \text{sc}, A)}$$

$$\frac{(X, \text{type}, A)}{(A, \text{sc}, A)}$$

Soundness and Completeness

Let \models denote the standard RDFS entailment, and $\vdash_{\rho\text{df}}$ a proof system based on the rules presented.

Theorem

Let G and H be graphs in ρdf then

$$G \models H \text{ if and only if } G \vdash_{\rho\text{df}} H.$$

Blank Nodes Modularization

Blank nodes can be treated in an orthogonal form to ρ df vocabulary.

Theorem

Let G and H be graphs in ρ df and $G \models H$, then

*there is a proof of H from G where the blank rule is used **at most once** and **at the end**.*

The role of reflexivity

The **only** consequence of reflexivity of sp and sc in RDFS is the possible entailment of triples of the form (x, sp, x) , (x, sc, x) .

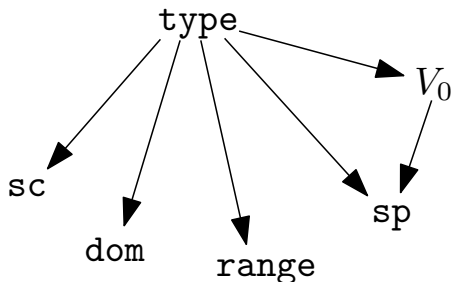
Theorem

Let G and H be ρ df graphs. Assume that H does not contain triples of the form (x, sp, x) and (x, sc, x) . Then,

$G \vdash_{\rho df} H$ without using reflexivity rules.

(Also, by not imposing reflexivity, axiomatic triples can be completely avoided.)

Dependence diagram among ρ df vocabulary



To determine $G \models H$ it is enough to test $G' \models H$ where G' is the subgraph of G which involves **only** nodes in $\text{voc}(H)$ and their dependencies in the diagram.

It is possible to avoid the closure to test RDFS entailment

- ▶ A naive approach to test $G \models H$ is:
 - ▶ (pre-)compute the closure of G
 - ▶ check if H is contained in the closure of G .

Theorem

The size of the closure of G is $O(|G|^2)$, and this bound is tight.

It is possible to avoid the closure to test RDFS entailment

- ▶ A naive approach to test $G \models H$ is:
 - ▶ (pre-)compute the closure of G
 - ▶ check if H is contained in the closure of G .

Theorem

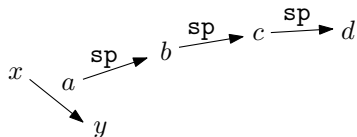
The size of the closure of G is $O(|G|^2)$, and this bound is tight.

- ▶ Alternative: to use a **goal oriented** approach based on the dependencies diagram.

Goal oriented entailment algorithm

Does the graph entails (x, d, y) ?

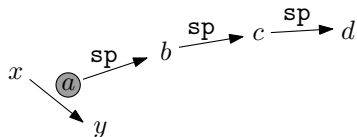
Look for triples of the form (x, a, y) and sp-paths from a to d .



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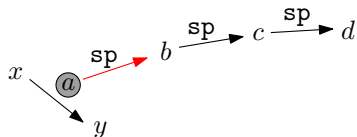
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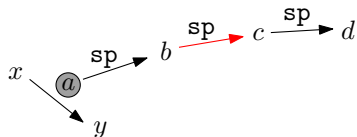
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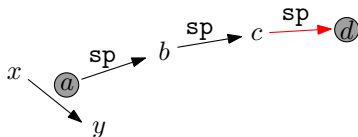
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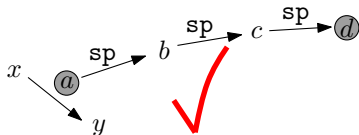
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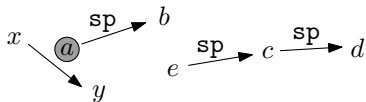
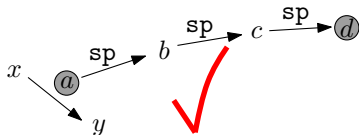
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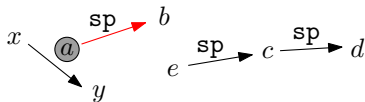
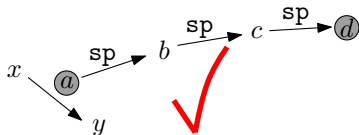
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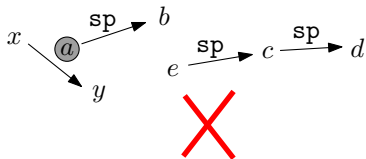
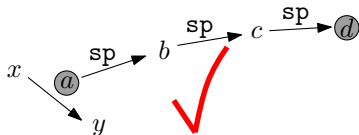
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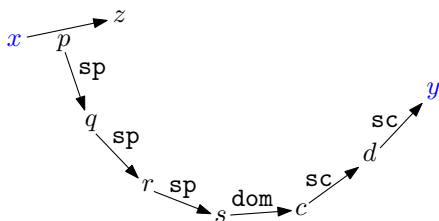
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Let $G = \{ \dots, (x, p, z), (p, sp, q), (q, sp, r), (r, sp, s), (s, dom, c), (c, sc, d), (d, sc, y), \dots \}$

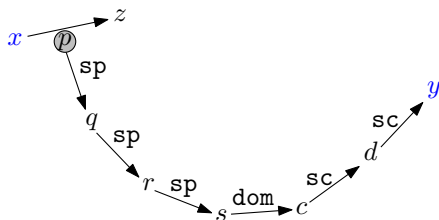
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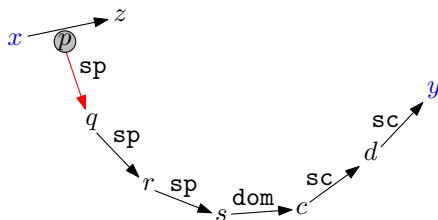
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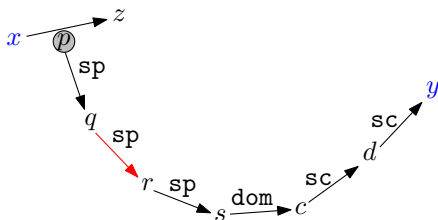
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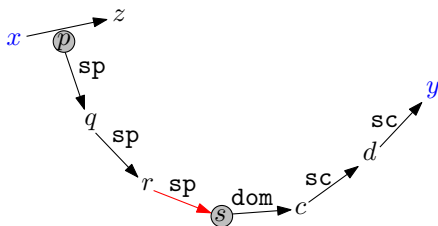
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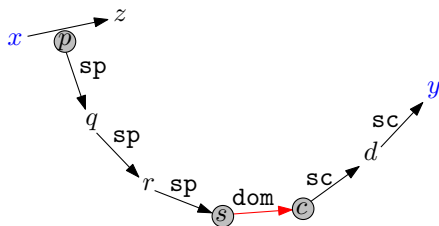
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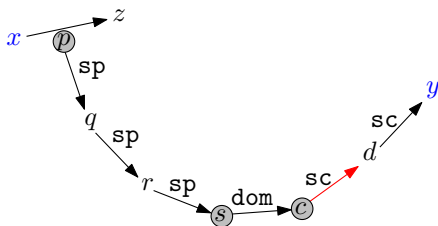
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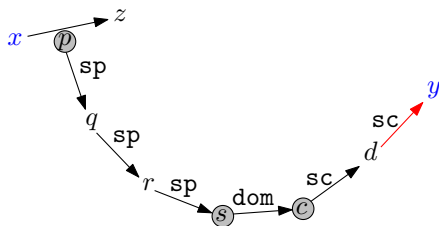
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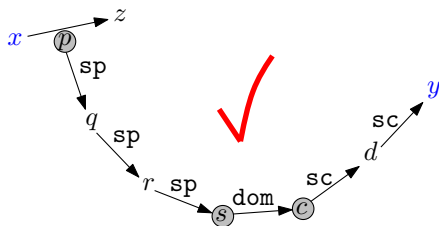
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Entailment can be done in $O(n \log n)$ time

Theorem

The goal oriented algorithm takes $O(|G| \log |G|)$ time in testing the entailment $G \models t$.

- ▶ Correctness follows by the dependencies diagram.
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- ▶ Time $|G|$ in traversing these data-structures.

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The $O(n \log n)$ bound is tight.

Theorem

Testing $G \models t$ takes time $\Omega(|G| \log |G|)$.

Idea: Coding the **set-disjointness** problem, which is $\Omega(n \log n)$

Conclusions

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 - ▶ Treat bnodes orthogonal to rest
- ▶ Next: Navigational language based on testing algorithm